

# Integrated Logic Systems (Part I)

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## Exercise 4.1

Download the constraint handling rules (CHRs) for  $\mathcal{ALC}$  Reasoning (`d1.chr`) from the course web page.

- a) Apply the CHRs by hand to show that the following statement leads to a contradiction:

```
isa(X,not(some(hasChild,human))), isa((X,Y),hasChild),
isa(Y,or(man,woman)).
```

- b) Use the CHRs to automatically show the truth of the following statements by contradiction:

- 1) Every proud parent has only happy children.
- 2) If Fritz is a parent and a child of Jody, then Jody is a grandparent.

## Exercise 4.2

For the following programs, determine all stable models by search over the truth-values of atoms and by using propagation. Show a complete search tree and all applied Propagation Rules.

- a)  $a :- b, \text{not } c.$   
 $b :- a, \text{not } c.$

- b)  $\text{dom}(1). \text{dom}(2).$   
 $p(X) :- \text{not } q(X), \text{dom}(X).$   
 $q(X) :- \text{not } p(X), \text{dom}(X).$   
 $:- p(1), p(2).$

- c)  $a :- \text{not } d.$   
 $a :- d, \text{not } e.$   
 $b :- e, \text{not } f.$   
 $c :- \text{not } b, \text{not } e.$   
 $c :- f, e, \text{not } a.$   
 $d :- \text{not } a, \text{not } c.$   
 $e :- \text{not } c.$   
 $f :- a, \text{not } b, \text{not } d.$

## Exercise 4.3

Let  $S$  be a stable model of some program  $P$ . Prove that  $S$  is a minimal Herbrand model of  $P$ .

*Hint:* Prove that

- 1)  $S$  is a Herbrand model of  $P$ , and that
- 2) there is no Herbrand model  $S'$  of  $P$  such that  $S' \subset S$  (by contradiction).