
Weighted Tree Automata

Exercise 31 (Theorem of Kleene)

1. Let $\Sigma = \{\sigma^{(2)}, \alpha^{(0)}, \beta^{(0)}\}$. Using the constructions from the lecture, show that $\{\sigma(\alpha, \beta)\}_{\beta}^* \cdot_{\beta} \{\alpha\}$ is in Rec.
2. Let $\Sigma = \{\sigma_1^{(2)}, \sigma_2^{(2)}, \alpha_1^{(0)}, \alpha_2^{(0)}\}$ and $G = (\{A_1, A_2\}, \Sigma, A_1, R)$ be a regular tree grammar with

$$R = \{A_1 \rightarrow \alpha_1, A_1 \rightarrow \sigma_1(A_1, A_2), A_2 \rightarrow \alpha_2, A_2 \rightarrow \sigma_2(A_2, A_1)\}.$$

Using the technique from the lecture, show that $L(G)$ is in Rat.

Exercise 32 (Myhill-Nerode theorem)

1. Let $\Sigma = \{\sigma^{(2)}, \alpha^{(0)}, \beta^{(0)}\}$ and $L \subseteq T_{\Sigma}$ be the language consisting of all trees with exactly as many α as β symbols. Use the Myhill-Nerode theorem to show that L is not recognizable.
2. But $\text{yield}(L)$ is a context-free language! Shouldn't L be recognizable by Thatcher's theorem $\text{yield}(\text{Rec}) = \text{CF}$?

Exercise 33 (Monadic second-order logic)

1. Let $\Sigma = \{\sigma^{(2)}, \gamma^{(1)}, \alpha^{(0)}, \beta^{(0)}\}$ and $\mathcal{V} = \{x, y, X\}$. Construct MSO-formulas $\varphi_1, \varphi_2, \varphi_3$ s.t. for every $\xi \in T_{\Sigma}$ and \mathcal{V} -assignment ρ for ξ ,
 - $(\xi, \rho) \models \varphi_1$ iff there is a downward path from the node $\rho(x)$ to the node $\rho(y)$ in ξ , i.e. there is a $w \in \mathbb{N}^*$ such that $\rho(y) = \rho(x)w$,
 - $(\xi, \rho) \models \varphi_2$ iff $\rho(X)$ is the set of all positions w in ξ such that $\xi|_w = \sigma(\alpha, \beta)$,
 - $(\xi, \rho) \models \varphi_3$ iff for every node in ξ labeled by σ , none of its child nodes is labeled by γ ,and $\text{Free}(\varphi_i) \subseteq \mathcal{V}$ for $i \in \{1, 2, 3\}$.
2. Consider the bu-deterministic fta $\mathcal{A} = (Q, \Sigma, \delta, F)$ with $Q = \{0, 1\}$, $F = \{1\}$, $\Sigma = \{\sigma^{(2)}, \gamma^{(1)}, \alpha^{(0)}\}$, and $\delta_{\sigma}(q_1, q_2) = \max\{q_1, q_2\}$ for every $q_1, q_2 \in Q$, $\delta_{\gamma}(q) = 1$ for every $q \in Q$, and $\delta_{\alpha}() = 0$.
Use the construction from the lecture to show that $L(\mathcal{A})$ is MSO-definable.
3. Let $\Sigma = \{\sigma^{(2)}, \gamma^{(1)}, \alpha^{(0)}\}$ and $\varphi = \exists x.\text{label}_{\gamma}(x)$. Use the construction from the lecture to show that $L(\varphi) \in \text{Rec}(\Sigma)$.